

# MATH 201 CLASS DISCUSSION: INDUCTION, CONTINUED

## 27 FEBRUARY 2020

### I Review exercises

- A) Prove that  $\forall n \geq 0 \quad 3|(n^3 - n)$
- B) Prove that  $\forall n \geq 0 \quad 2^{3n+1} + 5$  is always a multiple of 7.
- C) Consider the following definition of a recursive sequence:

$$\text{Let } a_1 = \frac{5}{2}.$$
$$\text{For } n \geq 1, \text{ let } a_{n+1} = \frac{1}{5}(a_n^2 + 6)$$

Prove that the sequence  $\{a_n\}$  is decreasing.

### II Find the flaw(s) in each of the following “proofs.”

- A) If any of  $n$  spiders is a tarantula, then all  $n$  spiders are tarantulas?
- B) I can lift all the sand on the beach.

**Proof.** Here we use the method of induction.

The proof is by induction.

For  $n \geq 1$  let  $P(n)$  be the predicate, “I can lift  $n$  grains of sand.”

**Base Case:**  $P(1)$  is true because I can certainly lift one grain of sand.

**Inductive Step:** Assume that I can lift  $n$  grains of sand.

I want to prove that I can lift  $n + 1$  grains of sand.

If I can lift  $n$  grains of sand, then surely I can lift  $n + 1$ ; one grain of sand will not make any difference.

Therefore  $P(n) \Rightarrow P(n + 1)$ .

By induction,  $P(n)$  is true for all  $n \geq 1$ .

- C) **Claim:** Given a set of  $n$  points in the plane, then these points are collinear (that is, lie on one line).

**Proof:** Here we use the method of induction.

For any non-negative integer,  $n$ , let  $P(n)$  assert that given any  $n$  points in the plane, then these  $n$  points lie on one line.

**Base case:** Clearly, if  $n = 1$ , one point lies on one line. So  $P(1)$  has been verified.

**Inductive Step:** Let  $n$  be a fixed positive integer.

Now assume that, given any set of  $n$  points, then these points are collinear.

Now consider a set of  $n + 1$  points.

Consider any subset of  $n$  points. Then these lie on a line. Consider another subset of  $n$  points. Then these too lie on a line.

The intersection of these sets contain  $n - 1$  points. So these lines are clearly the same.

Hence  $P(n+1)$

**Flaw?**

- D) Find the flaw in the following example of Donald Knuth, the distinguished computer scientist.

**Claim:**

$$\underbrace{\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \dots}_{n \text{ terms}} = \frac{3}{2} - \frac{1}{n}$$

**Base case:** For  $n = 1$ , we have  $\frac{3}{2} - \frac{1}{1} = \frac{1}{2}$

**Inductive step:**

$$\begin{aligned} \left( \frac{1}{1 \cdot 2} + \dots + \frac{1}{(n-1) \cdot n} \right) + \frac{1}{n \cdot (n+1)} &= \frac{3}{2} - \frac{1}{n} + \frac{1}{n \cdot (n+1)} \\ &= \frac{3}{2} - \frac{1}{n} + \frac{1}{n} - \frac{1}{n+1} = \frac{3}{2} - \frac{1}{n+1} \end{aligned}$$

**Flaw:**

- E) Find the flaw:

*(Incorrect!) Proof:* Let  $P(n)$  be the statement “the sum of the first  $n$  powers of two is  $2^n - 1$ .” We will prove by induction that  $P(n)$  holds for all  $n \in \mathbb{N}$ , from which the theorem follows.

For the base case, we prove  $P(0)$ , that the sum of the first zero powers of two is  $2^0 - 1$ . Since the sum of zero numbers is 0 and  $2^0 - 1 = 0$ , this result is true.

For the inductive step, assume that  $P(k)$  is true, meaning that

$$2^0 + 2^1 + \dots + 2^{k-1} = 2^k - 1. \quad (1)$$

We will prove that  $P(k+1)$  is true, meaning that the sum of the first  $k+1$  powers of two is  $2^{k+1} - 1$ . To see this, note that

$$\begin{aligned} 2^0 + 2^1 + \dots + 2^{k-1} + 2^k &= 2^k - 1 + 2^k \\ 2^0 + 2^1 + \dots + 2^{k-1} + 2^k &= 2(2^k) - 1 \\ 2^0 + 2^1 + \dots + 2^{k-1} + 2^k &= 2^k + 2^k - 1 \\ 2^0 + 2^1 + \dots + 2^{k-1} &= 2^k - 1 \end{aligned}$$

We've arrived at statement (1), which we know is true. Therefore,  $P(k+1)$  is true, completing the induction. ■

### III Strong Induction

- A) State the Principle of Strong Induction  
 B) Making Change (MIT)

The country Inductia, whose unit of currency is the Strong, has coins worth 3Sg (3 Strongs) and 5Sg. Although the Inductians have some trouble making small change like 4Sg or 7Sg, it turns out that they can collect coins to make change for any number that is at least 8 Strongs. Prove this.

- C) Find the flaw:

**Claim:** For every non-negative integer  $n$ ,  $5n = 0$ .

**Proof:** Here we use the method of strong induction

**Base case:** Let  $n = 0$ . Then  $5n = 0$ . Hence  $P(0)$  is true.

**Inductive step:** Let  $n \geq 0$  be a fixed integer.

Assume that  $5j = 0$  for all non-negative integers  $j$  with  $0 \leq j < n$ .

Write  $n + 1 = i + j$  where  $i$  and  $j$  are non-negative numbers, each less than  $n + 1$ .

Using the induction hypothesis,  
 $5(k + 1) = 5(i + j) = 5i + 5j = 0 + 0 = 0$

**Flaw**

D) Find the flaw in the following argument:

**Claim:**  $\frac{d}{dx} x^n = 0$  for all  $n \geq 0$ .

**Base case:** ( $n = 0$ ):  $\frac{d}{dx} x^0 = \frac{d}{dx} 1 = 0$

**Inductive step:** Assume that  $\frac{d}{dx} x^k = 0$  for all  $k \leq n$ . Then by the product rule,

$$\frac{d}{dx} x^{n+1} = \frac{d}{dx} (x^n \cdot x^1) = x^n \frac{d}{dx} x^1 + \left( \frac{d}{dx} x^n \right) x^1 = x^n \cdot 0 + 0 \cdot x^1 = 0.$$

**Flaw:**

Proof. First, let us find the value of  $n$  for which we will prove the statement. (does the statement hold for  $n = 1, 2, 3, \dots$ ? Have you found  $n$  for which this is true?) Let  $P(n) : 2n < n!$  be the statement. We will show that it holds for all  $n \geq \dots$  • **BASE STEP** (show  $P(n)$  holds for  $n$  smallest possible) • **INDUCTION HYPOTHESIS**  $P(k) : \dots$  (state the assumption for  $P(k)$ ) • **INDUCTION STEP** (keep in mind what you are trying to prove - it helps to note it on the side) (hint: notice that  $2 < k + 1 \forall k > 1$ ) • **CONCLUSION** (finish the proof by writing the conclusion)

Result true, proof false Stanford

In practice, it can be easy to inadvertently get this backwards. Here's an incorrect proof that the sum of the first  $n$  powers of two is  $2n - 1$ . (Note that the result that it proves is true, but the proof itself has a logical error that we'll discuss in a second).

(Incorrect!) Proof: Let  $P(n)$  be the statement "the sum of the first  $n$  powers of two is  $2n - 1$ ." We will prove by induction that  $P(n)$  holds for all  $n \in \mathbb{N}$ , from which the theorem follows. For the base case, we prove  $P(0)$ , that the sum of the first zero powers of two is  $2 \cdot 0 - 1$ . Since the sum of zero numbers is 0 and  $2 \cdot 0 - 1 = 0$ , this result is true. For the inductive step, assume that  $P(k)$  is true, meaning that  $2^0 + 2^1 + \dots + 2^{k-1} = 2k - 1$ . (1) We will prove that  $P(k+1)$  is true, meaning that the sum of the first  $k+1$  powers of two is  $2(k+1) - 1$ . To see this, note that  $2^0 + 2^1 + \dots + 2^{k-1} + 2^k = 2(k+1) - 1$ .  $2^0 + 2^1 + \dots + 2^{k-1} + 2^k = 2(2k) - 1$ .  $2^0 + 2^1 + \dots + 2^{k-1} + 2^k = 2k + 2k - 1$ .  $2^0 + 2^1 + \dots + 2^{k-1} + 2^k = 2k - 1$ . We've arrived at statement (1), which we know is true. Therefore,  $P(k+1)$  is true, completing the induction. ■ This proof is, unfortunately, incorrect, but it might not immediately be clear why. The setup of the

❖ Use Principle of Mathematical Induction to show that  $\left( \bigcup_{k=1}^n A_k \right)' = \bigcap_{k=1}^n A_k' \forall n \geq 2$ .

What is the principle of strong induction?

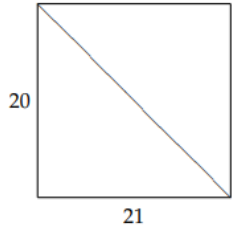
- ❖ Every integer greater than 1 is a product of primes
- ❖ Making Change (MIT)
- ❖ The country Inductia, whose unit of currency is the Strong, has coins worth 3Sg (3 Strongs) and 5Sg. Although the Inductians have some trouble making small change like 4Sg or 7Sg, it turns out that they can collect coins to make change for any number that is at least 8 Strongs.

Strong induction makes this easy to prove for  $n \in \mathbb{N}$ , because then  $n \in \mathbb{N} / 3 \geq 8$ , so by strong induction the Inductians can make change for exactly  $n \in \mathbb{N} / 3$  Strongs, and then they can add a 3Sg coin to get  $n \in \mathbb{N} / 3 + 1$  Sg. So the only thing to do is check that they can make change for all the amounts from 8 to 10Sg, which is not too hard to do. Here's a detailed writeup using the official form

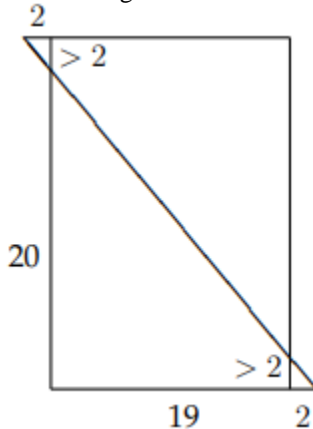
In principle, a proof should establish the truth of a proposition with absolute certainty. In practice, however, many purported proofs contain errors: overlooked cases, logical slips, and even algebra mistakes. But in a well-written proof, even if there is a bug, one should at least be able to pinpoint a specific statement that does not logically follow. See if you can find the first error in the following argument.

**MIT False Theorem 1.**  $420 > 422$  Proof. We will demonstrate this fact geometrically. We begin with a  $20 \times 21$  rectangle, which has area 420:

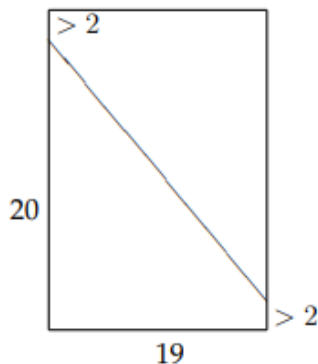
Proof. We will demonstrate this fact geometrically. We begin with a  $20 \times 21$  rectangle, which has area 420:



Now we cut along the diagonal as indicated above and slide the upper piece parallel to the cut until it has moved exactly 2 units leftward. This leaves a couple of stray corners, which are 2 units wide and just over 2 units high.



Finally, we snip off the two corners and place them together to form an additional small rectangle:



Now we have two rectangles, a large one with area just over  $(20 + 2) \times 19 = 418$  and a small one with area just over  $2 \times 2 = 4$ . Thus, the total area of the resulting figure is a bit over  $418 + 4 = 422$ . By conservation of area, 420 is equal to just a little bit more than 422. Where is the error?

Here is one published by Knuth.

**Claim:**

$$\underbrace{\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \dots}_{n \text{ terms}} = \frac{3}{2} - \frac{1}{n}$$

**Base case:** For  $n = 1$ , we have  $\frac{3}{2} - \frac{1}{1} = \frac{1}{1 \cdot 2}$

**Inductive step:**

$$\begin{aligned} \left( \frac{1}{1 \cdot 2} + \dots + \frac{1}{(n-1) \cdot n} \right) + \frac{1}{n \cdot (n+1)} &= \frac{3}{2} - \frac{1}{n} + \frac{1}{n \cdot (n+1)} \\ &= \frac{3}{2} - \frac{1}{n} + \frac{1}{n} - \frac{1}{n+1} = \frac{3}{2} - \frac{1}{n+1} \end{aligned}$$

**Flaw:**