

CLASS DISCUSSION

27 January 2020

Rates of growth: little o and big O notation

Suppose that $f(x) \rightarrow \infty$ and $g(x) \rightarrow \infty$ as $x \rightarrow \infty$. We say that “ f is of smaller order than g ” if

$\frac{f(x)}{g(x)} \rightarrow 0$ as $x \rightarrow \infty$. In this case, we write $f = o(g)$.

Now assume that f and g are each positive for large x . We say that “ f is at most the order of g ” if there is a positive integer M for which $\frac{f(x)}{g(x)} \leq M$ for large x . In this case, we write $f = O(g)$.

Little o and big O notation are common in computer science, as well as in discrete mathematics. Here is an example from the CS Dept of S.F. State:

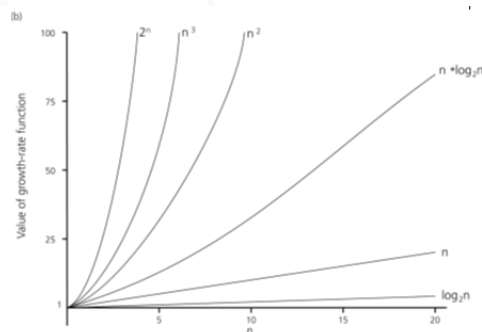
- $O(1)$ – **constant** time, the time is independent of n , e.g. array look-up
 - $O(\log n)$ – **logarithmic** time, usually the log is base 2, e.g. binary search
 - $O(n)$ – **linear** time, e.g. linear search
 - $O(n \log n)$ – e.g. efficient sorting algorithms
 - $O(n^2)$ – **quadratic** time, e.g. selection sort
 - $O(n^k)$ – **polynomial** (where k is some constant)
 - $O(2^n)$ – **exponential** time, very slow!
- Order of growth of some common functions
 $O(1) < O(\log n) < O(n) < O(n \log n) < O(n^2) < O(n^3) < O(2^n)$

➤ Order-of-Magnitude Analysis and Big O Notation

(a)

Function	n					
	10	100	1,000	10,000	100,000	1,000,000
1	1	1	1	1	1	1
$\log_2 n$	3	6	9	13	16	19
n	10	10^2	10^3	10^4	10^5	10^6
$n \cdot \log_2 n$	30	664	9,965	10^5	10^6	10^7
n^2	10^2	10^4	10^6	10^8	10^{10}	10^{12}
n^3	10^3	10^6	10^9	10^{12}	10^{15}	10^{18}
2^n	10^3	10^{30}	10^{301}	$10^{3,010}$	$10^{30,103}$	$10^{301,030}$

A comparison of growth-rate functions: a) in tabular form



A comparison of growth-rate functions: b) in graphical form

➤ Arithmetic of Big- O Notation

- 1) If $f(n)$ is $O(g(n))$ then $c \cdot f(n)$ is $O(g(n))$, where c is a constant.
 - Example: $23 \cdot \log n$ is $O(\log n)$
- 2) If $f_1(n)$ is $O(g(n))$ and $f_2(n)$ is $O(g(n))$ then also $f_1(n) + f_2(n)$ is $O(g(n))$
 - Example: what is order of $n^2 + n$?
 n^2 is $O(n^2)$
 n is $O(n)$ but also $O(n^2)$
 therefore $n^2 + n$ is $O(n^2)$

➤ **EXERCISES:**

Determine which of the following statements are true; justify each answer.

- (a) $3x^2 + 11 = o(x^5 + x + 99)$ (b) $x + 5 \sin x = O(x)$ (c) $2^x = o(x^{100})$
- (d) $3^x = O(e^x)$ (e) $x = o(\ln x)$ (f) $3 + \ln x + \ln(\ln x) + \sqrt{x} = o\left(x^{\frac{2}{3}}\right)$
- (g) $\ln x = o(\sqrt{x})$ (h) $(x^2 + 1)^4 = O((2x + 1)^3 x^5)$ (i) $\frac{x^2 + 13x + 2009}{5x + 1789} = O(\sqrt{x^2 + 9})$
- (j) $\ln x = o(\ln(\ln x))$ (k) $\ln(x^{55} + x^{33} + x^{11} + 1) = O(\ln x)$ (l) $7n - 2 = O(n)$
- (m) $3n^3 + 20n^2 + 5 = O(n^3)$ (n) $3 \ln n + \ln(\ln n) = O(\ln n)$

➤ **Exercises (Purdue University)**

Arrange the following list of functions in ascending order of growth rate, i.e. if function $g(n)$ immediately follows $f(n)$ in your list then, it should be the case that $f(n) = O(g(n))$.

$$g_1(n) = 2^{\sqrt{\ln n}} \quad g_2(n) = 2^n \quad g_3(n) = n^{\frac{4}{3}} \quad g_4(n) = n(\ln n)^3$$

$$g_5(n) = n^{\ln n} \quad g_6(n) = 2^{2^n} \quad g_7(n) = 2^{n^2}$$



[Edmund Landau](#) (1877 – 1938) is known for his work in analytic number theory and the distribution of primes. He first introduced the *little oh* notation in 1909.

